Balanceability

Antoine Dailly

Joint work with Laura Eslava¹, Adriana Hansberg², Alexandre Talon and Denae Ventura².

 1 IIMAS - Instituto de Investigaciones en Matemáticas Aplicadas y en Sistemas 2 Instituto de Matemáticas, UNAM Juriquilla





Context: Ramsey Theory

Idea

Guarantee the existence of ordered substructures within large chaotic structures.

Ramsey's Theorem (for 2-colorings) (1930)

For any k, there is an integer R(k) such that, if $n \ge R(k)$, then, every 2-edge-coloring of K_n contains a monochromatic K_k .



Context: extremal graph theory

Idea

Find the minimum density guaranteeing a given property, and the densest graphs for which it does not hold.

Turán's Theorem (1941) If *G* of order *n* contains more than $(1 - \frac{1}{k})\frac{n^2}{2}$ edges, then, *G* contains a K_{k+1} . The extremal graph is the balanced complete *k*-partite graph of order *n*.

Notations for the rest of the talk

- We consider 2-edge-colorings of K_n : $E(K_n) = R \sqcup B$.
- ► We denote by ex(n, G) the maximum number of edges in a G-free graph of order n.

Notations for the rest of the talk

- We consider 2-edge-colorings of K_n : $E(K_n) = R \sqcup B$.
- ► We denote by ex(n, G) the maximum number of edges in a G-free graph of order n.

Goal: generalizing Ramsey's ideas and looking for unavoidable patterns other than monochromatic copies.

Definition

An (r,b)-copy of a graph G(V,E) (with r+b=|E|) is a copy of G with r edges in R and b edges in B.

Definition

An (r, b)-copy of a graph G(V, E) (with r + b = |E|) is a copy of G with r edges in R and b edges in B.

⇒ By Ramsey, if *n* is large enough, we always have a (0, |E|)-copy or an (|E|, 0)-copy of *G*.

Definition

An (r, b)-copy of a graph G(V, E) (with r + b = |E|) is a copy of G with r edges in R and b edges in B.

⇒ By Ramsey, if *n* is large enough, we always have a (0, |E|)-copy or an (|E|, 0)-copy of *G*.

We want to guarantee the existence of an (r, b)-copy of G (for a given r > 0).

Definition

An (r, b)-copy of a graph G(V, E) (with r + b = |E|) is a copy of G with r edges in R and b edges in B.

⇒ By Ramsey, if *n* is large enough, we always have a (0, |E|)-copy or an (|E|, 0)-copy of *G*.

We want to guarantee the existence of an (r, b)-copy of G (for a given r > 0).

 \Rightarrow Need for a given density of each color class.

Definition

An (r,b)-copy of a graph G(V,E) (with r+b=|E|) is a copy of G with r edges in R and b edges in B.

⇒ By Ramsey, if *n* is large enough, we always have a (0, |E|)-copy or an (|E|, 0)-copy of *G*.

We want to guarantee the existence of an (r, b)-copy of G (for a given r > 0).

 \Rightarrow Need for a given density of each color class.

r-tonality

If, for every *n* large enough, there exists k(n,r) such that every 2-edge-coloring $R \sqcup B$ of K_n verifying |R|, |B| > k(n,r) contains an (r, b)-copy of *G*, then *G* is *r*-tonal.

Balanced copyA balanced copy of G(V, E) is an (r, b)-copy of G with $r \in \left\{ \lfloor \frac{|E|}{2} \rfloor, \lceil \frac{|E|}{2} \rceil \right\}.$

Balanced copy

A balanced copy of G(V, E) is an (r, b)-copy of G with $r \in \left\{ \left\lfloor \frac{|E|}{2} \right\rfloor, \left\lceil \frac{|E|}{2} \right\rceil \right\}.$

Balanceability (Caro, Hansberg, Montejano, 2020)

Let bal(n, G) be the smallest integer, if it exists, such that every 2-edge-coloring $R \sqcup B$ of K_n verifying |R|, |B| > bal(n, G) contains a balanced copy of G.

Balanced copy

A balanced copy of G(V, E) is an (r, b)-copy of G with $r \in \left\{ \left\lfloor \frac{|E|}{2} \right\rfloor, \left\lceil \frac{|E|}{2} \right\rceil \right\}.$

Balanceability (Caro, Hansberg, Montejano, 2020)

Let bal(n, G) be the smallest integer, if it exists, such that every 2-edge-coloring $R \sqcup B$ of K_n verifying |R|, |B| > bal(n, G) contains a balanced copy of G. If there is an n_0 such that, for every $n \ge n_0$, bal(n, G) exists, then G is balanceable

Balanced copy

A balanced copy of G(V, E) is an (r, b)-copy of G with $r \in \left\{ \left\lfloor \frac{|E|}{2} \right\rfloor, \left\lceil \frac{|E|}{2} \right\rceil \right\}.$

Balanceability (Caro, Hansberg, Montejano, 2020)

Let bal(n, G) be the smallest integer, if it exists, such that every 2-edge-coloring $R \sqcup B$ of K_n verifying |R|, |B| > bal(n, G) contains a balanced copy of G. If there is an n_0 such that, for every $n \ge n_0$, bal(n, G) exists, then

G is balanceable and bal(n, G) is its balancing number.

Balanced copy

A balanced copy of G(V, E) is an (r, b)-copy of G with $r \in \left\{ \left\lfloor \frac{|E|}{2} \right\rfloor, \left\lceil \frac{|E|}{2} \right\rceil \right\}.$

Balanceability (Caro, Hansberg, Montejano, 2020)

Let bal(n, G) be the smallest integer, if it exists, such that every 2-edge-coloring $R \sqcup B$ of K_n verifying |R|, |B| > bal(n, G) contains a balanced copy of G.

If there is an n_0 such that, for every $n \ge n_0$, bal(n, G) exists, then G is balanceable and bal(n, G) is its balancing number.

Ramsey-type problem

Extremal-type problem

Characterization

Theorem (Caro, Hansberg, Montejano, 2020)

A graph is balanceable if and only if it has both:



Characterization

Theorem (Caro, Hansberg, Montejano, 2020)

A graph is balanceable if and only if it has both: 1. A cut crossed by half of its edges;



Characterization

Theorem (Caro, Hansberg, Montejano, 2020)

A graph is balanceable if and only if it has both:

- 1. A cut crossed by half of its edges;
- 2. An induced subgraph containing half of its edges.



G is balanceable \Rightarrow

It has to fit in those two specific colorings of K_n :



G is balanceable \Rightarrow

It has to fit in those two specific colorings of K_n :



G is balanceable \Rightarrow

It has to fit in those two specific colorings of K_n :



G is balanceable \Rightarrow

It has to fit in those two specific colorings of K_n :



Those two specific colorings of K_n can be balanced $(|R| = |B| = \frac{1}{2} \binom{n}{2})$ for an infinity of values of n.

Theorem (Caro, Hansberg, Montejano, 2020)

For every *t*, there exists $\phi(n,t) \in \mathcal{O}(n^{2-\frac{1}{m(t)}})$ such that, if *n* is large enough, then, every 2-edge-coloring of K_n verifying $|R|, |B| \ge \phi(n,t)$ contains either a type A or a type B copy of K_{2t} .

Also shown (with a bound of $\epsilon\binom{n}{2}$) by Cutler & Montágh (2008) and Fox & Sudakov (2008).



Theorem (Caro, Hansberg, Montejano, 2020)

For every *t*, there exists $\phi(n,t) \in \mathcal{O}(n^{2-\frac{1}{m(t)}})$ such that, if *n* is large enough, then, every 2-edge-coloring of K_n verifying $|R|, |B| \ge \phi(n,t)$ contains either a type A or a type B copy of K_{2t} .

Also shown (with a bound of $\epsilon\binom{n}{2}$) by Cutler & Montágh (2008) and Fox & Sudakov (2008).



 \Rightarrow Gives a subquadratic bound for bal(*n*, *G*)

Theorem (Caro, Hansberg, Montejano, 2020)



Theorem (Caro, Hansberg, Montejano, 2020)



Theorem (Caro, Hansberg, Montejano, 2020)



Theorem (Caro, Hansberg, Montejano, 2020)



Theorem (Caro, Hansberg, Montejano, 2020)



Theorem (Caro, Hansberg, Montejano, 2020)



Theorem (Caro, Hansberg, Montejano, 2020)



Theorem (Caro, Hansberg, Montejano, 2020)



Theorem (Caro, Hansberg, Montejano, 2020)



Theorem (Caro, Hansberg, Montejano, 2020)



Theorem (Caro, Hansberg, Montejano, 2020)



Previous results on balanceability

- ► Caro, Hansberg, Montejano (2019)
 - ► $bal(n, K_4) = n 1$ or n (depending on $n \mod 4$)
 - No other complete graph with an even number of edges is balanceable!
Previous results on balanceability

- ► Caro, Hansberg, Montejano (2019)
 - ▶ $bal(n, K_4) = n 1$ or n (depending on $n \mod 4$)
 - No other complete graph with an even number of edges is balanceable!
- ► Caro, Hansberg, Montejano (2020)
 - ► Trees are balanceable For $n \ge 4k$, bal $(n, T_k) \le (k-1)n$
 - For k even and $n \ge \max(3, \frac{k^2}{4} + 1)$, $\operatorname{bal}(n, K_{1,k}) = \operatorname{bal}(n, K_{1,k+1}) = \left(\frac{k-2}{2}\right)n - \frac{k^2}{8} + \frac{k}{4}$

For
$$n \ge \frac{9}{32}k^2 + \frac{1}{4}k + 1$$
,
 $bal(n, P_{4k}) = bal(n, P_{4k+1}) = (k-1)n - \frac{1}{2}(k^2 - k - \frac{1}{2})$
 $bal(n, P_{4k-2}) = bal(n, P_{4k-1}) = (k-1)n - \frac{1}{2}(k^2 - k)$
 $\land P_k$ is the path on k edges (sorry \textcircled{C})

Previous results on balanceability

- ► Caro, Hansberg, Montejano (2019)
 - ▶ $bal(n, K_4) = n 1$ or n (depending on $n \mod 4$)
 - No other complete graph with an even number of edges is balanceable!
- ► Caro, Hansberg, Montejano (2020)
 - ► Trees are balanceable For $n \ge 4k$, bal $(n, T_k) \le (k-1)n$
 - ► For k even and $n \ge \max(3, \frac{k^2}{4} + 1)$, bal $(n, K_{1,k}) = bal(n, K_{1,k+1}) = \left(\frac{k-2}{2}\right)n - \frac{k^2}{8} + \frac{k}{4}$
 - ► For $n \ge \frac{9}{32}k^2 + \frac{1}{4}k + 1$, $bal(n, P_{4k}) = bal(n, P_{4k+1}) = (k-1)n - \frac{1}{2}(k^2 - k - \frac{1}{2})$ $bal(n, P_{4k-2}) = bal(n, P_{4k-1}) = (k-1)n - \frac{1}{2}(k^2 - k)$ $\land P_k$ is the path on k edges (sorry S)
- ► Caro, Lauri, Zarb (2020)

Balancing numbers of the graphs with at most 4 edges

Theorem (D., Eslava, Hansberg, Ventura, 2020+)





- C_{4k+2} is not balanceable ;
- $C_{4k+\epsilon}$ is balanceable
- C_{4k} is balanceable

Theorem (D., Eslava, Hansberg, Ventura, 2020+)

- C_{4k+2} is not balanceable ;
- ► $C_{4k+\epsilon}$ is balanceable, and bal $(n, C_{4k+\epsilon}) = (k-1)n - \frac{1}{2}(k^2 - k - 1 - \epsilon)$;
- C_{4k} is balanceable

Theorem (D., Eslava, Hansberg, Ventura, 2020+)

- C_{4k+2} is not balanceable ;
- ► $C_{4k+\epsilon}$ is balanceable, and bal $(n, C_{4k+\epsilon}) = (k-1)n - \frac{1}{2}(k^2 - k - 1 - \epsilon)$;

•
$$C_{4k}$$
 is balanceable, and
 $(k-1)n - (k-1)^2 \le bal(n, C_{4k}) \le (k-1)n + 12k^2 + 3k.$



Proposition

The cycle C_{4k+2} is not balanceable.

Proof by contradiction

Proposition

The cycle C_{4k+2} is not balanceable.

Proof by contradiction

 C_{4k+2} has a cut containing half of its edges.



Proposition

The cycle C_{4k+2} is not balanceable.

Proof by contradiction

 C_{4k+2} has a cut containing half of its edges.



Proposition

Let k be a positive integer, n be an integer such that $n \ge \frac{9}{2}k^2 + \frac{13}{4}k + \frac{49}{32}$, et $\epsilon \in \{-1, 1\}$. bal $(n, C_{4k+\epsilon})$ = bal $(n, P_{4k+\epsilon-1})$ = $(k-1)n - \frac{1}{2}(k^2 - k - 1 - \epsilon)$

Proposition

Let k be a positive integer, n be an integer such that $n \ge \frac{9}{2}k^2 + \frac{13}{4}k + \frac{49}{32}$, et $\epsilon \in \{-1, 1\}$.

$$bal(n, C_{4k+\epsilon}) = bal(n, P_{4k+\epsilon-1})$$
$$= (k-1)n - \frac{1}{2}(k^2 - k - 1 - \epsilon)$$

Proof (for C_{4k+1})



Balanced $P_{4k} \Rightarrow$ 2k edges of each color

Proposition

Let k be a positive integer, n be an integer such that $n \ge \frac{9}{2}k^2 + \frac{13}{4}k + \frac{49}{32}$, et $\epsilon \in \{-1, 1\}$.

$$bal(n, C_{4k+\epsilon}) = bal(n, P_{4k+\epsilon-1})$$
$$= (k-1)n - \frac{1}{2}(k^2 - k - 1 - \epsilon)$$

Proof (for C_{4k+1})



Balanced $P_{4k} \Rightarrow$ 2k edges of each color

Proposition

Let k be a positive integer, n be an integer such that $n \ge \frac{9}{2}k^2 + \frac{13}{4}k + \frac{49}{32}$, et $\epsilon \in \{-1, 1\}$.

$$bal(n, C_{4k+\epsilon}) = bal(n, P_{4k+\epsilon-1})$$
$$= (k-1)n - \frac{1}{2}(k^2 - k - 1 - \epsilon)$$

Proof (for C_{4k+1})



Balanced $P_{4k} \Rightarrow$ 2k edges of each color

We can close the cycle which will be balanced

Proposition

Let k be a positive integer, n be an integer such that $n \ge \frac{9}{2}k^2 + \frac{13}{4}k + \frac{49}{32}$, et $\epsilon \in \{-1, 1\}$.

$$bal(n, C_{4k+\epsilon}) = bal(n, P_{4k+\epsilon-1})$$
$$= (k-1)n - \frac{1}{2}(k^2 - k - 1 - \epsilon)$$

Proof (for C_{4k+1})

Balanced $C_{4k+1} \Rightarrow$ A color with 2k+1 edges



Proposition

Let k be a positive integer, n be an integer such that $n \ge \frac{9}{2}k^2 + \frac{13}{4}k + \frac{49}{32}$, et $\epsilon \in \{-1, 1\}$.

$$bal(n, C_{4k+\epsilon}) = bal(n, P_{4k+\epsilon-1})$$
$$= (k-1)n - \frac{1}{2}(k^2 - k - 1 - \epsilon)$$

Proof (for C_{4k+1})

Balanced $C_{4k+1} \Rightarrow$ A color with 2k+1 edges



Proposition

Let k be a positive integer, n be an integer such that $n \ge \frac{9}{2}k^2 + \frac{13}{4}k + \frac{49}{32}$, et $\epsilon \in \{-1, 1\}$.

$$bal(n, C_{4k+\epsilon}) = bal(n, P_{4k+\epsilon-1})$$
$$= (k-1)n - \frac{1}{2}(k^2 - k - 1 - \epsilon)$$

Proof (for C_{4k+1})

 $\begin{array}{l} \text{Balanced } C_{4k+1} \Rightarrow \\ \text{A color with } 2k+1 \text{ edges} \\ \text{Removing one gives} \\ \text{a balanced } P_{4k} \end{array}$



The proof for odd cycles does not work:

The proof for odd cycles does not work:

Balanced P_{4k-1}



The proof for odd cycles does not work:

Balanced P_{4k-1} \Rightarrow The cycle may be non-balanced



The proof for odd cycles does not work:



Theorem (D., Eslava, Hansberg, Ventura, 2020+)
For
$$n \ge \frac{9}{2}k^2 + \frac{13}{4}k + \frac{49}{32}$$
,
 $(k-1)n - (k-1)^2 \le bal(n, C_{4k}) \le (k-1)n + 12k^2 + 3k$

Proposition

For every $n \ge 4k$, $bal(n, C_{4k}) \ge (k-1)n - (k-1)^2$.

Proposition

For every $n \ge 4k$, $bal(n, C_{4k}) \ge (k-1)n - (k-1)^2$.

Proof

We build a 2-edge-coloring $R \sqcup B$ with no balanced C_{4k} and such that $|B| \ge |R| = (k-1)n - (k-1)^2$.

Proposition

For every
$$n \ge 4k$$
, $bal(n, C_{4k}) \ge (k-1)n - (k-1)^2$.

Proof

We build a 2-edge-coloring $R \sqcup B$ with no balanced C_{4k} and such that $|B| \ge |R| = (k-1)n - (k-1)^2$.



Proposition

For every
$$n \ge 4k$$
, $bal(n, C_{4k}) \ge (k-1)n - (k-1)^2$.

Proof

We build a 2-edge-coloring $R \sqcup B$ with no balanced C_{4k} and such that $|B| \ge |R| = (k-1)n - (k-1)^2$.



 \Rightarrow A cycle can have at most 2k-2 edges in *R*.

Cycles C_{4k} : upper bound (1)



Proof by contradiction

Cycles C_{4k} : upper bound (1)

Proposition
Let
$$k > 0$$
 and $n \ge \frac{9}{2}k^2 + \frac{13}{4}k + \frac{49}{32}$:
bal $(n, C_{4k}) \le (k-1)n + 12k^2 + 3k$.

Proof by contradiction $|R|, |B| > bal(n, P_{4k-2}) \Rightarrow$ There is a balanced P_{4k-2} .

$$4k-1$$
 vertices
 $O-- O-O$

Cycles C_{4k} : upper bound (1)



Proof by contradiction

- $|R|, |B| > bal(n, P_{4k-2}) \Rightarrow$ There is a balanced P_{4k-2} .
 - \Rightarrow We close it with (wlog) a *B* edge

$$\underbrace{4k-1 \text{ vertices}}_{\mathbf{O}_{----}} \left. \begin{array}{c} 2k-1 \text{ in } R\\ 2k \text{ in } B \end{array} \right.$$







Lemmas enforce the colors of E(X), E(Y) and E(X, Y).



We cannot have $|X|, |Y| \ge k$



We cannot have $|X|, |Y| \ge k \Rightarrow$ wlog, assume |X| < k





Consider the graph induced by $(C_{4k+1} \cup X, Y) \cap \mathbb{R}$.



Consider the graph induced by $(C_{4k+1} \cup X, Y) \cap \mathbb{R}$.

It contains $\geq (k-1)n$ edges; and $\exp(n, P_{2k-1}) \leq (k-1)n$ [FS13]
Cycles C_{4k} : upper bound (2) Proof by contradiction (sequel)



Consider the graph induced by $(C_{4k+1} \cup X, Y) \cap \mathbb{R}$.

It contains $\geq (k-1)n$ edges; and $\exp(n, P_{2k-1}) \leq (k-1)n$ [FS13] \Rightarrow It contains a P_{2k-1} .

Cycles C_{4k} : upper bound (2) Proof by contradiction (sequel)



Consider the graph induced by $(C_{4k+1} \cup X, Y) \cap R$.

It contains $\ge (k-1)n$ edges; and $ex(n, P_{2k-1}) \le (k-1)n$ [FS13] \Rightarrow It contains a P_{2k-1} . There remain enough edges in R to have a $K_{1,2}$.

Cycles C_{4k} : upper bound (2) Proof by contradiction (sequel)



Consider the graph induced by $(C_{4k+1} \cup X, Y) \cap R$.

It contains $\geq (k-1)n$ edges; and $\exp(n, P_{2k-1}) \leq (k-1)n$ [FS13] \Rightarrow It contains a P_{2k-1} .

There remain enough edges in R to have a $K_{1,2}$.

We complete with edges in Y, which will be in B, and we get a balanced C_{4k} .

 \Rightarrow Contradiction

$\frac{\text{Definition}}{C_{k,\ell} \text{ is a cycle } C_k}$



Definition

 $C_{k,\ell}$ is a cycle C_k with the $u_i u_{i+\ell}$ chords.



Definition

 $C_{k,\ell}$ is a cycle C_k with the $u_i u_{i+\ell}$ chords.



Contains antiprisms and Möbius ladders.

Theorem (D., Hansberg, Ventura, 2021)

Let k > 3 and $\ell \in \{2, ..., k-2\}$. The graph $C_{k,\ell}$ is balanceable if and only if k is even and $(k, \ell) \neq (6, 2)$.

Theorem (D., Hansberg, Ventura, 2021)

Let k > 3 and $\ell \in \{2, ..., k-2\}$. The graph $C_{k,\ell}$ is balanceable if and only if k is even and $(k, \ell) \neq (6, 2)$.

Proof in eight cases! Each case uses the characterization.

Theorem (D., Hansberg, Ventura, 2021)

Let k > 3 and $\ell \in \{2, ..., k-2\}$. The graph $C_{k,\ell}$ is balanceable if and only if k is even and $(k, \ell) \neq (6, 2)$.

Proof in eight cases! Each case uses the characterization.

Theorem (D., Hansberg, Ventura, 2021)

Let k > 3 and $\ell \in \{2, ..., k-2\}$. The graph $C_{k,\ell}$ is balanceable if and only if k is even and $(k, \ell) \neq (6, 2)$.

Proof in eight cases! Each case uses the characterization.

Proof of the case k = 4a, ℓ even

Proposition

If, in G(V, E), I is an independent set such that $\sum_{u \in I} d(u) = \frac{|E|}{2}$, then, G is balanceable.

Theorem (D., Hansberg, Ventura, 2021)

Let k > 3 and $\ell \in \{2, ..., k-2\}$. The graph $C_{k,\ell}$ is balanceable if and only if k is even and $(k, \ell) \neq (6, 2)$.

Proof in eight cases! Each case uses the characterization.



Theorem (D., Hansberg, Ventura, 2021)

Let k > 3 and $\ell \in \{2, ..., k-2\}$. The graph $C_{k,\ell}$ is balanceable if and only if k is even and $(k, \ell) \neq (6, 2)$.

Proof in eight cases! Each case uses the characterization.



Theorem (D., Hansberg, Ventura, 2021)

Let k > 3 and $\ell \in \{2, ..., k-2\}$. The graph $C_{k,\ell}$ is balanceable if and only if k is even and $(k, \ell) \neq (6, 2)$.

Proof in eight cases! Each case uses the characterization.



► K_n with $\frac{n(n-1)}{2}$ odd Common integer solutions of $k(n-k) = \frac{1}{2} \binom{n}{2} \pm \frac{1}{2}$ and $\binom{\ell}{2} = \frac{1}{2} \binom{n}{2} \pm \frac{1}{2}$

- ► K_n with $\frac{n(n-1)}{2}$ odd Common integer solutions of $k(n-k) = \frac{1}{2} \binom{n}{2} \pm \frac{1}{2}$ and $\binom{\ell}{2} = \frac{1}{2} \binom{n}{2} \pm \frac{1}{2}$
 - \Rightarrow Explicit, but difficult to combine... found by computation:
 - 1. $n \in \{2, 3, 7, 11, 14, 38, 62, 79, 359, 43.262\} \Rightarrow$ Balanceable
 - 2. Other values of $n \le 10^{765.500} \Rightarrow$ Non-balanceable

- ► K_n with $\frac{n(n-1)}{2}$ odd Common integer solutions of $k(n-k) = \frac{1}{2} \binom{n}{2} \pm \frac{1}{2}$ and $\binom{\ell}{2} = \frac{1}{2} \binom{n}{2} \pm \frac{1}{2}$
 - \Rightarrow Explicit, but difficult to combine... found by computation:
 - 1. $n \in \{2, 3, 7, 11, 14, 38, 62, 79, 359, 43.262\} \Rightarrow$ Balanceable
 - 2. Other values of $n \le 10^{765.500} \Rightarrow$ Non-balanceable



- ► K_n with $\frac{n(n-1)}{2}$ odd Common integer solutions of $k(n-k) = \frac{1}{2} \binom{n}{2} \pm \frac{1}{2}$ and $\binom{\ell}{2} = \frac{1}{2} \binom{n}{2} \pm \frac{1}{2}$
 - \Rightarrow Explicit, but difficult to combine... found by computation:
 - 1. $n \in \{2, 3, 7, 11, 14, 38, 62, 79, 359, 43.262\} \Rightarrow$ Balanceable
 - 2. Other values of $n \le 10^{765.500} \Rightarrow$ Non-balanceable
- ► 2K_n

Balanceable \Leftrightarrow *n* is the sum of two squares

- K_n with n(n-1)/2 odd Common integer solutions of k(n-k) = 1/2 (n/2) ± 1/2 and (^ℓ/₂) = 1/2 (n/2) ± 1/2 ⇒ Explicit, but difficult to combine... found by computation:
 - 1. $n \in \{2, 3, 7, 11, 14, 38, 62, 79, 359, 43, 262\} \Rightarrow Balanceable$
 - 2. Other values of $n \le 10^{765.500} \Rightarrow$ Non-balanceable

► 2K_n

Balanceable \Leftrightarrow *n* is the sum of two squares

 \Rightarrow Allows us to break graph operators (disjoint union, joint...)!

Summary



Summary



Summary



Non-balanceable graphs

Theorem (Caro, Hansberg, Montejano, 2020)

A graph is balanceable if and only if it has both:

- C_{4k+2} has the induced subgraph, not the cut
- ► K₅ has neither

Non-balanceable graphs

Theorem (Caro, Hansberg, Montejano, 2020)

A graph is balanceable if and only if it has both: 1. A cut crossed by half of its edges;

- C_{4k+2} has the induced subgraph, not the cut
- ► K₅ has neither

 \rightarrow Can we differentiate "levels" of non-balanceability?

Non-balanceable graphs

Theorem (Caro, Hansberg, Montejano, 2020)

A graph is balanceable if and only if it has both:

1. A cut crossed by half of its edges;

2. An induced subgraph containing half of its edges.

- C_{4k+2} has the induced subgraph, not the cut
- ► K₅ has neither

 \rightarrow Can we differentiate "levels" of non-balanceability?

Idea

From a 2-edge-coloring to a 2-edge-covering:

Idea

From a 2-edge-coloring to a 2-edge-covering:

1. The edges are labeled with $\{r\}$, $\{b\}$ or $\{r, b\}$.

Idea

From a 2-edge-coloring to a 2-edge-covering:

- 1. The edges are labeled with $\{r\}$, $\{b\}$ or $\{r, b\}$.
- 2. Edges labeled with {*r*, *b*} are called bicolored; we can choose their color.

Idea

From a 2-edge-coloring to a 2-edge-covering:

- 1. The edges are labeled with $\{r\}$, $\{b\}$ or $\{r, b\}$.
- 2. Edges labeled with {*r*, *b*} are called bicolored; we can choose their color.
- 3. Denoting the 2-edge-covering by $c, R = \{e \mid r \in c(e)\}$ and $B = \{e \mid b \in c(e)\}.$

Idea

From a 2-edge-coloring to a 2-edge-covering:

- 1. The edges are labeled with $\{r\}$, $\{b\}$ or $\{r, b\}$.
- 2. Edges labeled with {*r*, *b*} are called bicolored; we can choose their color.
- 3. Denoting the 2-edge-covering by $c, R = \{e \mid r \in c(e)\}$ and $B = \{e \mid b \in c(e)\}.$

Definition (D., Eslava, Hansberg, Ventura, 2020+)

Let $bal^*(n, G)$ be the smallest integer such that every 2-edgecovering $R \cup B$ of K_n verifying $|R|, |B| > bal^*(n, G)$ contains a balanced copy of G. $bal^*(n, G)$ is called the generalized balancing number of G.

Idea

From a 2-edge-coloring to a 2-edge-covering:

- 1. The edges are labeled with $\{r\}$, $\{b\}$ or $\{r, b\}$.
- 2. Edges labeled with {*r*, *b*} are called bicolored; we can choose their color.
- 3. Denoting the 2-edge-covering by $c, R = \{e \mid r \in c(e)\}$ and $B = \{e \mid b \in c(e)\}.$

Definition (D., Eslava, Hansberg, Ventura, 2020+)

Let $bal^*(n, G)$ be the smallest integer such that every 2-edgecovering $R \cup B$ of K_n verifying $|R|, |B| > bal^*(n, G)$ contains a balanced copy of G.

 $bal^*(n, G)$ is called the generalized balancing number of G.

 \Rightarrow Every graph has a generalized balancing number!

Proposition

If bal(n, G) exists, then $bal^*(n, G) = bal(n, G)$. Otherwise, $\frac{1}{2} \binom{n}{2} \le bal^*(n, G) < \binom{n}{2}$.

Proposition

If bal(n, G) exists, then $bal^*(n, G) = bal(n, G)$. Otherwise, $\frac{1}{2} \binom{n}{2} \le bal^*(n, G) < \binom{n}{2}$.

Counting bicolored edges If $|R|, |B| = \frac{1}{2} \binom{n}{2} + b$:

$$\begin{array}{c} R \setminus B & -b \\ \frac{1}{2} \binom{n}{2} - b & B \setminus R \\ b & -\frac{1}{2} \binom{n}{2} - b \end{array}$$

Proposition

If bal(n, G) exists, then $bal^*(n, G) = bal(n, G)$. Otherwise, $\frac{1}{2} \binom{n}{2} \le bal^*(n, G) < \binom{n}{2}$.

Counting bicolored edges If $|R|, |B| = \frac{1}{2} \binom{n}{2} + b$:



 $\Rightarrow 2b$ bicolored edges

Proposition

If bal(n, G) exists, then $bal^*(n, G) = bal(n, G)$. Otherwise, $\frac{1}{2} \binom{n}{2} \le bal^*(n, G) < \binom{n}{2}$.

Counting bicolored edges If $|R|, |B| = \frac{1}{2} \binom{n}{2} + b$:

$$\begin{array}{c|c} R \setminus B & b \\ \hline \frac{1}{2} \binom{n}{2} - b & B \setminus R \\ b & -\frac{1}{2} \binom{n}{2} - b \end{array} \Rightarrow \begin{array}{c} 2b \text{ bicolored} \\ edges \end{array}$$

Proposition

If k bicolored edges guarantee a balanced copy of G, then $bal^*(n, G) \leq \frac{1}{2} \binom{n}{2} + \left\lceil \frac{k}{2} \right\rceil - 1.$

A general upper bound

►
$$\mathcal{H}(G) = \left\{ H \le G \mid e(H) = \left\lfloor \frac{e(G)}{2} \right\rfloor, H \text{ with no isolated vertex} \right\}$$
A general upper bound

►
$$\mathcal{H}(G) = \left\{ H \le G \mid e(H) = \left\lfloor \frac{e(G)}{2} \right\rfloor, H \text{ with no isolated vertex} \right\}$$

Theorem (D., Eslava, Hansberg, Ventura, 2020+)

For every G(V, E) and $n \ge |V|$, we have

$$\mathsf{bal}^*(n,G) \leq \frac{1}{2} \binom{n}{2} + \left\lceil \frac{\mathsf{ex}(n,\mathscr{H}(G))}{2} \right\rceil.$$

A general upper bound

►
$$\mathcal{H}(G) = \left\{ H \le G \mid e(H) = \left\lfloor \frac{e(G)}{2} \right\rfloor, H \text{ with no isolated vertex} \right\}$$

Theorem (D., Eslava, Hansberg, Ventura, 2020+)

For every G(V, E) and $n \ge |V|$, we have

$$\mathsf{bal}^*(n,G) \leq \frac{1}{2} \binom{n}{2} + \left\lceil \frac{\mathsf{ex}(n,\mathcal{H}(G))}{2} \right\rceil$$

Proof

If there are at least $ex(n, \mathcal{H}(G)) + 1$ bicolored edges, we can select them, complete the copy of G, and assign the colors of the bicolored edges to balance the copy.

Cycles C_{4k+2}

• $\mathcal{H}(C_{4k+2}) = \text{ linear forests of size } 2k+1$

Cycles C_{4k+2}

- $\mathcal{H}(C_{4k+2}) = \text{ linear forests of size } 2k+1$
- ▶ bal* $(n, C_{4k+2}) \le \frac{1}{2} \binom{n}{2} + \frac{kn}{2} + \mathcal{O}(1)$ [Ning, Wang, 2020]

Cycles C_{4k+2}

- $\mathcal{H}(C_{4k+2}) = \text{ linear forests of size } 2k+1$
- ► bal* $(n, C_{4k+2}) \le \frac{1}{2} \binom{n}{2} + \frac{kn}{2} + \mathcal{O}(1)$ [Ning, Wang, 2020]

Cycles C_{4k+2}

- $\mathcal{H}(C_{4k+2}) =$ linear forests of size 2k+1
- ► bal* $(n, C_{4k+2}) \le \frac{1}{2} \binom{n}{2} + \frac{kn}{2} + \mathcal{O}(1)$ [Ning, Wang, 2020]



Cycles C_{4k+2}

- $\mathcal{H}(C_{4k+2}) =$ linear forests of size 2k+1
- ► bal* $(n, C_{4k+2}) \le \frac{1}{2} \binom{n}{2} + \frac{kn}{2} + \mathcal{O}(1)$ [Ning, Wang, 2020]

► bal*
$$(n, \mathcal{K}_5) \leq \frac{1}{2} \binom{n}{2} + (1 + \epsilon) \frac{1}{4\sqrt{2}} n^{\frac{3}{2}}$$
 [Füredi, Simonovits, 2013]

Cycles C_{4k+2}

- $\mathcal{H}(C_{4k+2}) =$ linear forests of size 2k+1
- ► bal* $(n, C_{4k+2}) \le \frac{1}{2} \binom{n}{2} + \frac{kn}{2} + \mathcal{O}(1)$ [Ning, Wang, 2020]

► bal*
$$(n, K_5) \leq \frac{1}{2} {n \choose 2} + (1 + \epsilon) \frac{1}{4\sqrt{2}} n^{\frac{3}{2}}$$
 [Füredi, Simonovits, 2013]

 \rightarrow Quality of this bound?

Theorem (D., Eslava, Hansberg, Ventura, 2020+)
bal*
$$(n, C_{4k+2}) = \frac{1}{2} \binom{n}{2}$$

Theorem (D., Eslava, Hansberg, Ventura, 2020+)
bal*
$$(n, C_{4k+2}) = \frac{1}{2} \binom{n}{2}$$

Proof idea

1. Assume $|R|, |B| > \frac{1}{2} \binom{n}{2}$: there are at least 2 bicolored edges.

Theorem (D., Eslava, Hansberg, Ventura, 2020+)
bal*
$$(n, C_{4k+2}) = \frac{1}{2} \binom{n}{2}$$

Proof idea

- 1. Assume $|R|, |B| > \frac{1}{2} \binom{n}{2}$: there are at least 2 bicolored edges.
- 2. Let $t \ge 3k + 1$. We find a type A or type B copy of K_{2t} .

Theorem (D., Eslava, Hansberg, Ventura, 2020+)
bal*
$$(n, C_{4k+2}) = \frac{1}{2} \binom{n}{2}$$

Proof idea

- 1. Assume $|R|, |B| > \frac{1}{2} \binom{n}{2}$: there are at least 2 bicolored edges.
- 2. Let $t \ge 3k + 1$. We find a type A or type B copy of K_{2t} .

3. If type A copy: balanced copy of C_{4k+2} .

Theorem (D., Eslava, Hansberg, Ventura, 2020+)
bal*
$$(n, C_{4k+2}) = \frac{1}{2} \binom{n}{2}$$

Proof idea

- 1. Assume $|R|, |B| > \frac{1}{2} \binom{n}{2}$: there are at least 2 bicolored edges.
- 2. Let $t \ge 3k + 1$. We find a type A or type B copy of K_{2t} .

- 3. If type A copy: balanced copy of C_{4k+2} .
- If type B copy: wherever the bicolored edge is, we can find a balanced copy of C_{4k+2}.

Proposition
Let
$$c = \left(\frac{\sqrt{2}-1}{2\sqrt{2}}\right)^{\frac{5}{2}}$$
. We have bal* $(n, K_5) \ge \frac{1}{2}{n \choose 2} + (1-\epsilon)cn^{\frac{3}{2}}$.

Proposition
Let
$$c = \left(\frac{\sqrt{2}-1}{2\sqrt{2}}\right)^{\frac{5}{2}}$$
. We have $bal^*(n, K_5) \ge \frac{1}{2}\binom{n}{2} + (1-\epsilon)cn^{\frac{3}{2}}$.

Proof

We build a 2-edge-covering of K_n with no balanced K_5 .

Proposition
Let
$$c = \left(\frac{\sqrt{2}-1}{2\sqrt{2}}\right)^{\frac{5}{2}}$$
. We have $bal^*(n, K_5) \ge \frac{1}{2}{n \choose 2} + (1-\epsilon)cn^{\frac{3}{2}}$.

Proof

We build a 2-edge-covering of K_n with no balanced K_5 .



Proposition
Let
$$c = \left(\frac{\sqrt{2}-1}{2\sqrt{2}}\right)^{\frac{5}{2}}$$
. We have $bal^*(n, K_5) \ge \frac{1}{2}\binom{n}{2} + (1-\epsilon)cn^{\frac{3}{2}}$.

Proof

We build a 2-edge-covering of K_n with no balanced K_5 .



Proposition
Let
$$c = \left(\frac{\sqrt{2}-1}{2\sqrt{2}}\right)^{\frac{5}{2}}$$
. We have $bal^*(n, K_5) \ge \frac{1}{2}\binom{n}{2} + (1-\epsilon)cn^{\frac{3}{2}}$.

Proof

We build a 2-edge-covering of K_n with no balanced K_5 .



Then, we prove $|R|, |B| > \frac{1}{2} \binom{n}{2} + (1-\epsilon) c n^{\frac{3}{2}}$.

Cycles C_{4k+2}

- General upper bound: bal* $(n, C_{4k+2}) \leq \frac{1}{2} {n \choose 2} + \frac{kn}{2} + \mathcal{O}(1)$
- Exact value: $bal^*(n, C_{4k+2}) = \frac{1}{2} \binom{n}{2}$

Cycles C_{4k+2}

- General upper bound: bal* $(n, C_{4k+2}) \le \frac{1}{2} \binom{n}{2} + \frac{kn}{2} + \mathcal{O}(1)$
- Exact value: $bal^*(n, C_{4k+2}) = \frac{1}{2} \binom{n}{2}$

K_5

- ► General upper bound: bal* $(n, K_5) \le \frac{1}{2} \binom{n}{2} + (1+\epsilon) \frac{1}{4\sqrt{2}} n^{\frac{3}{2}}$
- Lower bound: bal* $(n, K_5) \ge \frac{1}{2} \binom{n}{2} + (1 \epsilon) \left(\frac{\sqrt{2} 1}{2\sqrt{2}}\right)^{\frac{5}{2}} n^{\frac{3}{2}}$

Cycles C_{4k+2}

- General upper bound: bal* $(n, C_{4k+2}) \leq \frac{1}{2} {n \choose 2} + \frac{kn}{2} + \mathcal{O}(1)$
- Exact value: $bal^*(n, C_{4k+2}) = \frac{1}{2} \binom{n}{2}$

K_5

- ► General upper bound: bal* $(n, K_5) \le \frac{1}{2} \binom{n}{2} + (1+\epsilon) \frac{1}{4\sqrt{2}} n^{\frac{3}{2}}$
- Lower bound: bal* $(n, K_5) \ge \frac{1}{2} {n \choose 2} + (1 \epsilon) \left(\frac{\sqrt{2} 1}{2\sqrt{2}}\right)^{\frac{5}{2}} n^{\frac{3}{2}}$

Cycles C_{4k+2}

- General upper bound: bal* $(n, C_{4k+2}) \leq \frac{1}{2} {n \choose 2} + \frac{kn}{2} + \mathcal{O}(1)$
- Exact value: $bal^*(n, C_{4k+2}) = \frac{1}{2} \binom{n}{2}$

K_5

- General upper bound: bal* $(n, K_5) \leq \frac{1}{2} {n \choose 2} + (1+\epsilon) \frac{1}{4\sqrt{2}} n^{\frac{3}{2}}$
- Lower bound: bal* $(n, K_5) \ge \frac{1}{2} \binom{n}{2} + (1 \epsilon) \left(\frac{\sqrt{2} 1}{2\sqrt{2}}\right)^{\frac{5}{2}} n^{\frac{3}{2}}$

 \rightarrow There are differences among non-balanceable graphs.

Final words

Conclusion

- ► Balanceability results, study of bal(*n*, *G*)
- Introduction of $bal^*(n, G)$ to study non-balanceable graphs

Final words

Conclusion

- ► Balanceability results, study of bal(*n*, *G*)
- Introduction of $bal^*(n, G)$ to study non-balanceable graphs

Open questions

- ► Complexity
- More graph classes
- More colors

Final words

Conclusion

- ► Balanceability results, study of bal(*n*, *G*)
- Introduction of $bal^*(n, G)$ to study non-balanceable graphs

Open questions

- ► Complexity
- More graph classes
- More colors

